

**IN THE SPECIFICATION**

Please replace paragraph no. 0004 with as follows:

U.S. Application No. 10/354,797, Methods and Systems of Host Caching, filed on January 29, 2003, now U.S. Patent No. 6,965,979 B2;

Please replace paragraph no. 0008 with as follows:

U.S. Applcation No. 10/616,128, Snapshots of File Systems in Data Storage Systems, filed on July 8, 2003, now U.S. Patent No. 6,959,313 B2; and

Please replace paragraph no. 0066 with as follows:

If in write back cache mode, the secondary host [1] 2 writes the block to the nonvolatile cache lines 280 (Figure 7) of the secondary host 2 at step 64. At step 66, the secondary host 2 marks the block dirty in the nonvolatile cache lines 280 of the secondary host 1 to indicate the block can be destaged to the secondary target VLUN.

Please replace paragraph no. 0069 with as follows:

Figure 10 illustrates further processing of a sync command identified at step 52 (Figure 8). At step 84, the secondary host 1 reads the data set ID of the sync command and checks if it is new at step 86. If it is [a] new, the secondary host 1 checks for availability of a new data structure at step 98. If a data structure is not available, the secondary host 1 rejects the sync command at step 104, discards the sync command at step 106 and awaits the next command at step 50 (Figure 8). If a new data structure is available at step 98, the secondary host 1 allocates the new data structure and stores the data set ID of the sync command in the data structure at step 102. Regardless of whether or not the data set ID is new, the secondary host 1 reads the sequence number of the sync command at step 85 and sets the corresponding identifier bitmap bit position from 0 to 1 at step 87. At step 72, the secondary host 1 manages the data sets as described above and shown in Figure 9. At step 73, the secondary host 1 acknowledges to the primary host that the processing of the sync command is complete and returns to step 50 (Figure 8) to wait for the next command.